NFA with  $\in$  moves is exactly same as NFA without  $\in$  moves.

But differece exist in the transition function  $\delta$ .  $\delta$  must include information about  $\in$  transitions.

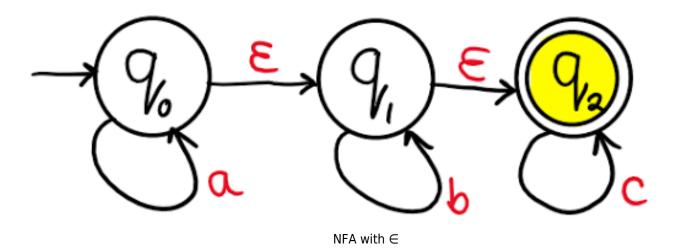
NFA with  $\in$ -Moves has 6 tuples  $(Q, \Sigma, \delta, q0, F)$ .

Where,

- Q = finite set of states.
- $\Sigma$  = finite set input symbols.
- $\delta$  = transition function that maps  $Q \times (\Sigma \cup \{\in\})$  to  $2^{\circ}$ .
- q0 = initial state.
- F = set of final states.

The non-deterministic finite automaton can be extended to include the transitions on  $\frac{1}{2}$  null/empty input  $\in$ .

For example,



In this NFA with epsilon,

- It accept an input string 'aabc'.
- Or string as number of a's followed by number of b's followed by number of c's.
- The string 'aabc' is accepted by the NFA by following the path with labels a, a, ∈, b, ∈,
  c.

Transition table for above NFA.

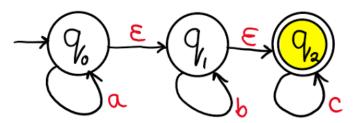


### ∈-closure

 $\in$ -closure of a state q is a set of states following by all transitions of q that are labeled as  $\in$ .

- $\in$ -closure (q0) = (q0, q1, q2)
- $\in$ -closure (q1) = (q1, q2)
- $\in$ -closure (q2) = (q2)

# NFA with $\in$ to NFA without $\in$



NFA with  $\in$ 

Transition diagram

## Transition table NFA with $\in$

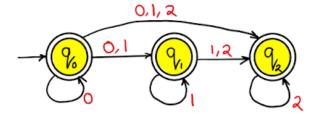


First find out ∈ closure:∈-closure

- (q0) = (q0, q1, q2)
- $\in$ -closure (q1) = (q1, q2)
- $\in$ -closure (q2) = (q2)

### Transition table NFA without ∈

State	а	b	С
<del>&gt;</del> @	{q0, q1, q2}	{q1, q2}	{q2}
<b>(1)</b>	Ф	{q1, q2}	{q2}
<u>@</u>	ф	ф	{q2}



NFA without ∈

#### Transition diagram

#### Reference:

• Introduction to the Theory of Computation" by Michael Sipser.

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